

# Introduction to QAOA

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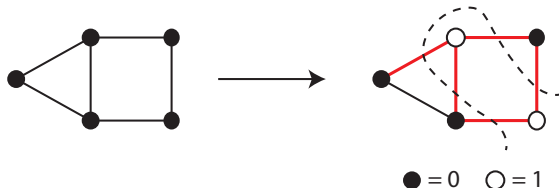
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## (Weighted) MaxCut

- Graph: Set of *vertices* ( $V$ ) connected by (weighted) *edges* ( $E$ )
- MaxCut: Partition of vertices into two disjoint subsets (labeled by 0 and 1), such that the total *weight* of the edges between the two subsets is maximized
- For equal-weight edges, the goal is simply to maximize the number of edges connecting the two subsets.



**Input:**

(Weighted) graph  $G = (V, E)$

**Output:**

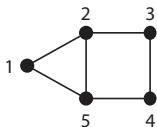
Maximum cut  $x \in (0, 1)^n$   
 $x = [0, 1, 0, 1, 0]$

- Quadratic Programs: Optimize (maximize or minimize) a quadratic objective function subject to linear constraints on the variables  
e.g., minimize  $\mathbf{x}^T Q \mathbf{x} + \mathbf{c}^T \mathbf{x}$   
subject to  $A \mathbf{x} \leq \mathbf{b}$
- Special case: Quadratic Unconstrained Binary Optimization (QUBO)
  - ▶ Quadratic objective function
  - ▶ No constraints on variables
  - ▶ Binary variables
- QUBO examples
  - ▶ MaxCut
  - ▶ Number partitioning
  - ▶ Graph coloring
  - ▶ ...

# MaxCut as QUBO

MaxCut

Weight matrix:



$$W = \begin{pmatrix} 0 & 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 \\ 1 & 1 & 0 & 1 & 0 \end{pmatrix}$$

Cost function:

$$C(\mathbf{x}) = \sum_{i,j=1}^n W_{ij} x_i (1 - x_j)$$

QUBO

QUBO matrix and vectors:

$$c_i = \sum_{j=1}^n W_{ij}, \quad Q_{ij} = -W_{ij}$$

Cost function:

$$\begin{aligned} C(\mathbf{x}) &= \mathbf{x}^T Q \mathbf{x} + \mathbf{c}^T \mathbf{x} \\ &= \sum_{i,j=1}^n Q_{ij} x_i x_j + \sum_{i=1}^n c_i x_i \end{aligned}$$

- Quantum Approximate Optimization Algorithm (QAOA) first introduced in Farhi, Goldstone and Gutmann (2014) [1]
- Finds approximate solutions to QUBO instances (e.g. MaxCut)
- Can be regarded as a special case of Variational Quantum Eigensolvers (VQE)
- Layerized variational form based on Trotterized adiabatic process (related to adiabatic quantum computing)
- Key idea: encode the cost function of the optimization problem in the cost Hamiltonian  $H_C$

## QUBO to Hamiltonian

Goal: Find the cost Hamiltonian operator  $H_C$  that encodes the cost function  $C(x)$ , i.e.,

$$H_C |x\rangle = C(x) |x\rangle, \quad x = \{0, 1\}.$$

Using the fact that

$$Z_i |x\rangle = (-1)^{x_i} |x\rangle = (1 - 2x_i) |x\rangle \implies \frac{1 - Z_i}{2} |x\rangle = x_i |x\rangle,$$

we have

$$\begin{aligned} C(x) &= \sum_{i,j} Q_{ij} x_i x_j + \sum_i c_i x_i \\ \implies H_C &= \sum_{i,j} Q_{ij} \left( \frac{1 - Z_i}{2} \right) \left( \frac{1 - Z_j}{2} \right) + \sum_i c_i \left( \frac{1 - Z_i}{2} \right) \\ &= \sum_{i,j} \frac{1}{4} Q_{ij} Z_i Z_j - \sum_i \frac{1}{2} \left( c_i + \sum_j Q_{ij} \right) Z_i + \left( \sum_{i,j} \frac{Q_{ij}}{4} + \sum_i \frac{c_i}{2} \right). \end{aligned}$$

## Trotterized AQC

Typical steps of Trotterized adiabatic quantum computing (AQC): [2]

- Prepare initial state as the highest energy state of some “mixer” Hamiltonian (that does not commute with  $H_C$ ),  $H_M = \sum_i X_i$ , i.e.,  $|\psi_0\rangle = \otimes_i |+\rangle$ .
- Set the total Hamiltonian  $H(t) = f(t)H_C + g(t)H_M$  with slowly varying control functions  $f(t) = t/T$  and  $g(t) = 1 - t/T$ .
- Through *adiabatic evolution*, the system will end up in the highest energy state of the cost Hamiltonian  $H_C$ , which then solves the QUBO problem.

# Trotterized AQC

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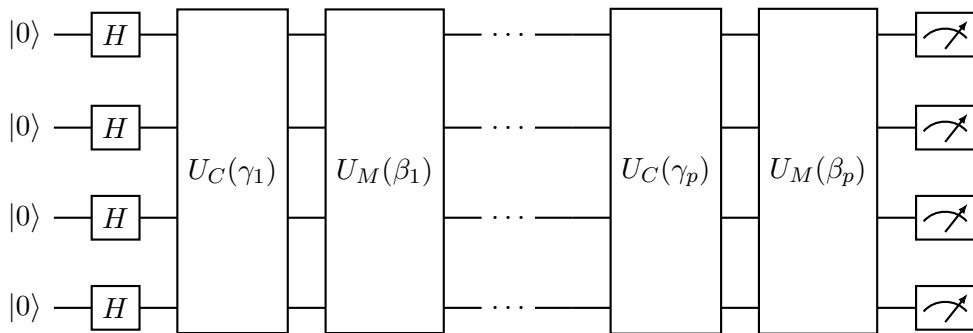
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- Through *adiabatic evolution*, the system will end up in the highest energy state of the cost Hamiltonian  $H_C$ , which then solves the QUBO problem.
- In practice, to implement the adiabatic evolution, one decomposes the time-evolution operator (for a time-dependent Hamiltonian) into a sequence of small steps through the *Trotter-Suzuki formula*:

$$\begin{aligned} U(t) := \mathcal{T} \exp \left[ -i \int_0^T H(t) dt \right] &\approx \prod_{a=0}^{k-1} \exp [-iH(a\tau)\tau] \\ &= \prod_{a=0}^{k-1} \exp [-if(a\tau)H_C\tau] \exp [-ig(a\tau)H_M\tau]. \end{aligned}$$



## QAOA Variational Form

Inspired by Trotterized AQC, QAOA was designed to be a variational algorithm with repeated cost and mixer layers.



$p$  repetitions of alternating cost and mixer layers:

$$U_C(\gamma_i) = e^{-i\gamma_i H_C}$$

$$U_M(\beta_i) = e^{-i\beta_i H_M}$$

# Matrix Exponentiation

Recall that

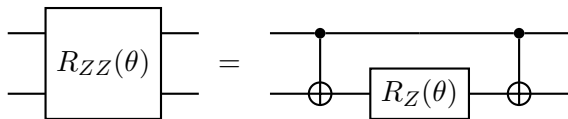
$$H_M = \sum_{i=1}^n X_i$$

$$H_C = \sum_{i,j=1}^n \frac{1}{4} Q_{ij} Z_i Z_j - \sum_{i=1}^n \frac{1}{2} \left( c_i + \sum_{j=1}^n Q_{ij} \right) Z_i + \left( \sum_{i,j=1}^n \frac{Q_{ij}}{4} + \sum_{i=1}^n \frac{c_i}{2} \right)$$

Then upon matrix exponentiation, the mixer and cost layers become

$$U_M(\beta) = e^{-i\beta H_M} = \prod_{i=1}^n R_{X_i}(2\beta)$$



$$U_C(\gamma) = e^{-i\gamma H_C} = \prod_{i,j=1}^n R_{Z_i Z_j} \left( \frac{1}{2} Q_{ij} \gamma \right) \prod_{i=1}^n R_{Z_i} \left( - \left( c_i + \sum_{j=1}^n Q_{ij} \right) \gamma \right).$$



## QAOA Workflow

- 1 Initialize  $\beta$  and  $\gamma$  with suitable real values.
- 2 Prepare the state  $|\psi(\beta, \gamma)\rangle$  using the QAOA circuit and measure it in the computational basis.
- 3 Compute the expectation value  $\langle\psi(\beta, \gamma)|H_C|\psi(\beta, \gamma)\rangle$ .
- 4 Find a new set of parameters  $(\beta_{\text{new}}, \gamma_{\text{new}})$  with a classical optimization algorithm (e.g., gradient descent) and use this new set of parameters in the QAOA circuit.
- 5 Repeat steps 2 - 4 until some suitable convergence criterion is met.
- 6 The solution is then approximated as  $|\psi(\beta_{\text{opt}}, \gamma_{\text{opt}})\rangle$  which maximizes  $\langle\psi(\beta_{\text{opt}}, \gamma_{\text{opt}})|H_C|\psi(\beta_{\text{opt}}, \gamma_{\text{opt}})\rangle$

## References

-  E. Farhi, J. Goldstone, and S. Gutmann, *A Quantum Approximate Optimization Algorithm*, arXiv:1411.4028 [quant-ph] (2014).
-  Y. Sun et al., *Adiabatic Quantum Simulation Using Trotterization*, arXiv:1805.11568 [quant-ph] (2018).

Materials presented also borrow from:

- Qiskit Global Summer School 2021, *Introduction to the Quantum Approximate Optimization Algorithm and Applications*.
- Learn Quantum Computation using Qiskit, Ch. 4.1.3, *Solving combinatorial optimization problems using QAOA*.